

---

## First Problem Set Assignment: BNF

---

---

### Learning Abstract

This assignment features defining BNF and its functions. It also demonstrates making a parse tree using the BNF created.

---

### Problem 1: Shapes

---

---

---

#### BNF Description

---

$\langle \text{SL} \rangle ::= \langle \text{empty} \rangle \mid ( \langle \text{style} \rangle )$

$\langle \text{style} \rangle ::= ( \text{size } \langle \text{size} \rangle ) \langle \text{style} \rangle \mid ( \text{color } \langle \text{color} \rangle ) \langle \text{style} \rangle \mid ( \text{pattern } \langle \text{pattern} \rangle ) \langle \text{style} \rangle \mid ( \text{shape } \langle \text{shape} \rangle ) \langle \text{style} \rangle$

$\langle \text{style} \rangle ::= \langle \text{empty} \rangle$

$\langle \text{size} \rangle ::= \text{large} \mid \text{medium} \mid \text{small}$

$\langle \text{color} \rangle ::= \text{red} \mid \text{blue} \mid \text{yellow}$

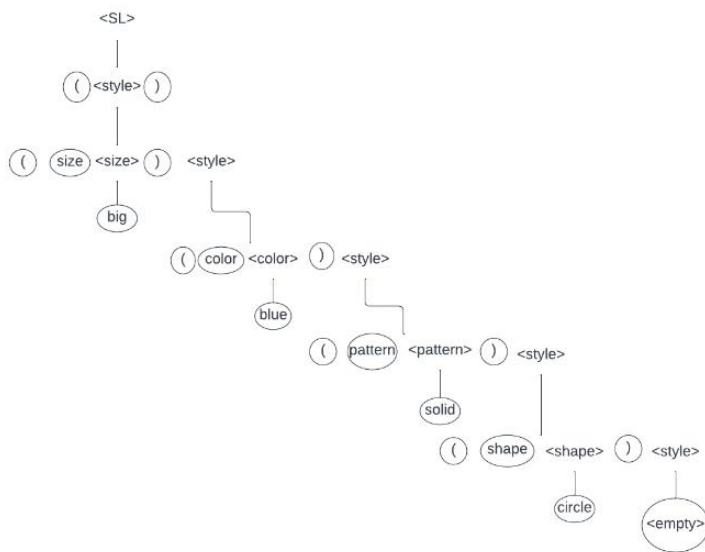
$\langle \text{pattern} \rangle ::= \text{stripped} \mid \text{dotted} \mid \text{solid}$

$\langle \text{shape} \rangle ::= \text{circle} \mid \text{triangle} \mid \text{square}$

---

### Parse Tree for Shapes (1)

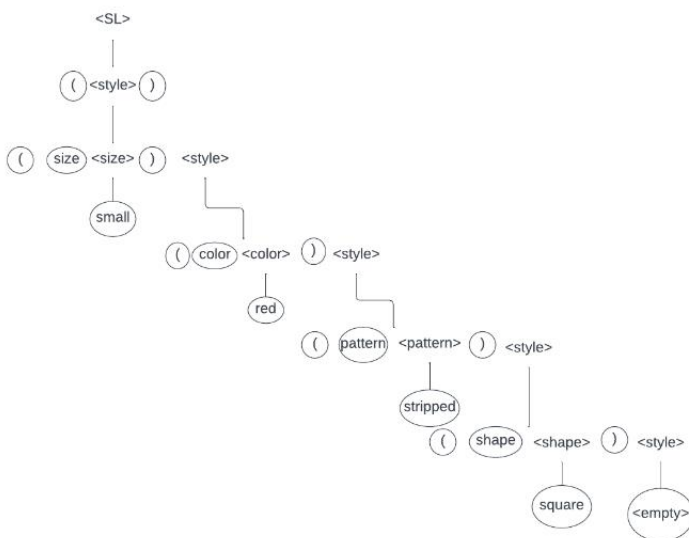
---




---

## Parse Tree for Shapes (2)

---




---

## Problem 2 - SQN (Special Quaternary Numbers)

---



---

### BNF Description

---

`<SQN> ::= <empty> | <D> | 0 <D0> | 1 <D1> | 2 <D2> | 3 <D3>`

`<D> ::= 0 | 1 | 2 | 3`

`<D0> ::= 1 <D1> | 2 <D2> | 3 <D3> | <empty>`

`<D1> ::= 0 <D0> | 2 <D2> | 3 <D3> | <empty>`

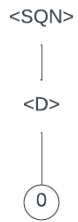
`<D2> ::= 0 <D0> | 1 <D1> | 3 <D3> | <empty>`

`<D3> ::= 0 <D0> | 1 <D1> | 2 <D2> | <empty>`

---

### Parse Tree for SQN (1)

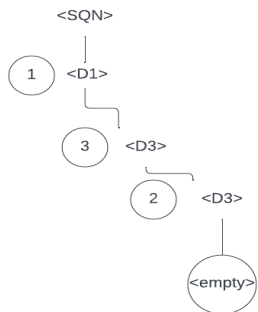
---



---

### Parse Tree for SQN (2)

---



---

### Parse Tree for SQN (3)

---

The string “1223” cannot be drawn because the production rule for <D2> won’t allow another 2 to come after.

---

## Problem 3: Fours

---

---

### BNF Description

---

$\langle S \rangle ::= \langle \text{empty} \rangle \mid \langle L1 \rangle \langle L2 \rangle \langle L3 \rangle \langle L4 \rangle$

$\langle L1 \rangle ::= (1\ 1\ 1\ 1) \mid \langle \text{empty} \rangle$

$\langle L2 \rangle ::= (1\ 1\ 2) \langle L2 \rangle \mid (1\ 2\ 1) \langle L2 \rangle \mid (2\ 1\ 1) \langle L2 \rangle \mid \langle \text{empty} \rangle$

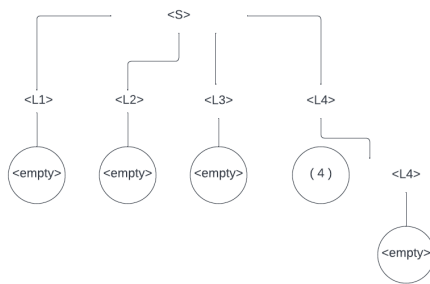
$\langle L3 \rangle ::= (3\ 1) \langle L3 \rangle \mid (1\ 3) \langle L3 \rangle \mid \langle \text{empty} \rangle$

$\langle L4 \rangle ::= (4) \langle L4 \rangle \mid \langle \text{empty} \rangle$

---

### Parse Tree for Fours (1)

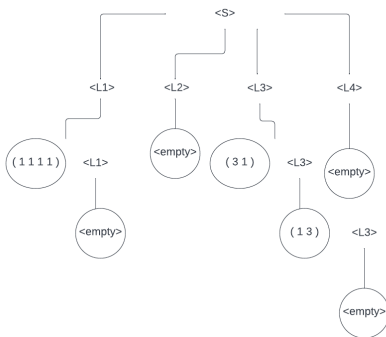
---



---

### Parse Tree for Fours (2)

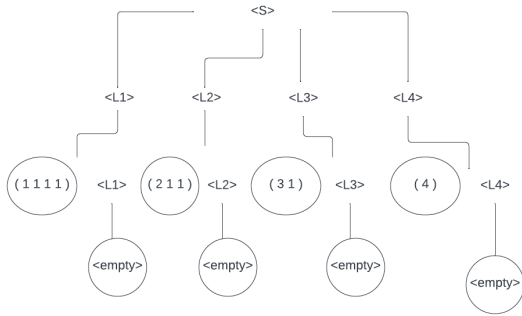
---



---

## Parse Tree for Fours (3)

---



---

## Problem 4: BXR

---

---

### BNF Description

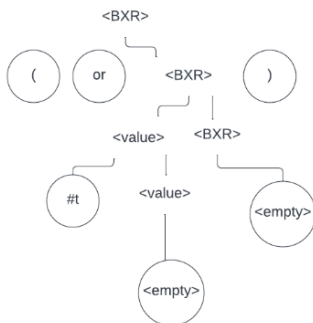
---

$\langle \text{BXR} \rangle ::= ( \text{ and } \langle \text{BXR} \rangle ) \mid ( \text{ or } \langle \text{BXR} \rangle ) \mid ( \text{ not } \#f ) \langle \text{BXR} \rangle \mid ( \text{ not } \#t ) \langle \text{BXR} \rangle \mid \langle \text{value} \rangle \langle \text{BXR} \rangle \mid \langle \text{empty} \rangle$   
 $\langle \text{value} \rangle ::= \#t \langle \text{value} \rangle \mid \#f \langle \text{value} \rangle \mid \langle \text{empty} \rangle$

---

### Parse Tree for BXR (1)

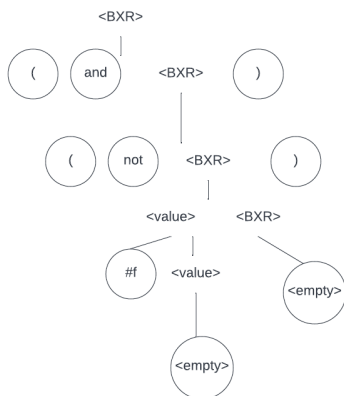
---



---

### Parse Tree for BXR (2)

---




---

## Problem 5: CF (Color Fun)

---



---

### BNF Description

---

<CF> ::= <add> | <describe> | <show> | colors | exit  
 <add> ::= add ( <color\_RGB> <color\_RGB> <color\_RGB> ) <color\_name> | add color <color\_name>  
 <describe> ::= describe <color\_name>  
 <show> ::= show <color\_name>  
 <color\_RGB> ::= 0 | 1 | 2 | 3 | ... | 256

---

### Parse Tree for CF (1)

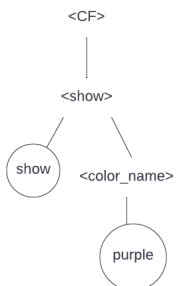
---




---

### Parse Tree for CF (2)

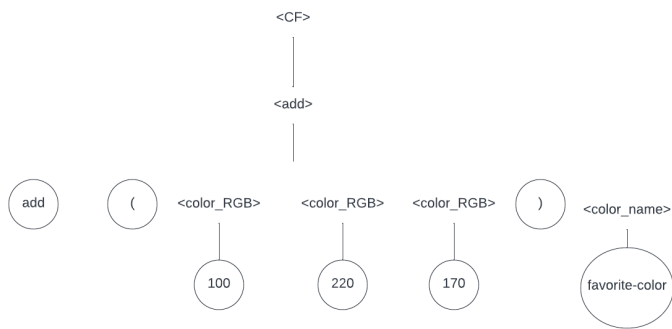
---




---

### Parse Tree for CF (3)

---



---

## Problem 6: BNF?

---

---

### BNF Definition

---

BNF, short for Backus–Naur Form, was created by a program designer named John Bakus. It's the formal technique used in computer science for structuring and describing the grammar syntax of any programming language. It contains sets of terminal symbols, nonterminal symbols, tokens, and the production rule. The production rule gets used to represent the left-hand nonterminal side being replaced by what is on the right-hand terminal or nonterminal side. After that, a parse tree gets constructed that breaks down the long expression into tiny parts based on the BNF made for it.